### **Continued Fractions in Lean**

A Newbie's Adventure

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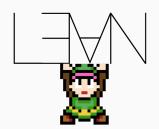
Let's Go on an Adventure

















...perhaps because I am interning at VU Amsterdam

• Some experience using Isabelle

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- · First project with a dependent type theorem prover

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- Basic maths and functional programming knowledge



# **Definitions**

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- each  $b_i$  is a partial denominator

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$$292 + \frac{1}{1 + \cdots}$$

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2 variable (a : Type\*)

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```
/- Fix a type -/
variable (a : Type*)

/-- A gcf_pair consists of a partial numerator a

→ and partial denominator b -/
structure gcf_pair := (a : a) (b : a)
```

```
-- Once a sequence hits none, it stays none def seq := {f : \mathbb{N} \to \text{option } \alpha \text{ // } \forall \text{ } \{n\}, \text{ } f \text{ } n = \text{ } none \to f \text{ } (n+1) = \text{ } none}
```

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\rightarrow f \text{ (n + 1) = none}

\rightarrow \text{ A generalized continued fraction consists of a}

\rightarrow \text{ leading head term (the "integer part") and a}

\rightarrow \text{ sequence of partial partial numerators } a_n \text{ and}

\rightarrow \text{ partial denominators } b_n \text{ -/}

\rightarrow \text{ structure gcf := (head : } \alpha) \text{ (seq : seq (gcf_pair } \rightarrow \alpha))}
```

#### **Evaluate Generalized Continued Fractions**

```
1 def convergents (g : gcf \alpha) (n : \mathbb{N}) : \alpha := 2 g.head + if n = 0 then 0 else aux n g.seq
```

#### **Evaluate Generalized Continued Fractions**

```
1 def aux : \mathbb{N} → seq (qcf_pair \alpha) → \alpha
2 | 0 s := match s.head with
3
  I none := 0
    | some \langle a, b \rangle := a / b
    end
6 | (n + 1) s := match s.head with
7 | none := 0
    | some \langle a, b \rangle := a / (b + aux n s.tail)
    end
9
10
II def convergents (g : gcf a) (n : \mathbb{N}) : a :=
12 g.head + if n = 0 then 0 else aux n g.seq
```

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2 def cf := {g : gcf \alpha // \forall (n : \mathbb{N}) (a : \alpha), \rightarrow (partial_numerators g).nth n = some a \rightarrow a = 1}
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First impression:

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First impression: Pretty Sweet!

## So, since cf is a subtype of gcf, we can do

```
1 def convergents (g : gcf a) (n : \mathbb{N}) : a := ...
2 
3 variable (c : cf a)
4 #check convergents c 0
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#### Oh, I see - I need to cast!

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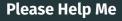
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...alright, let's go on Zulip 🗷



A few minutes and messages from Kevin Buzzard later...

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```
instance cf_to_gcf : has_coe (cf β) (gcf β)
2 := by {unfold cf, apply_instance}
3
4 /- Best practice: create a lemma for your cast -/
5 @[simp, elim_cast]
6 lemma coe_cf (c : cf β) : (†c : gcf β) = c.val
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#### Now this works:

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#### This, however, still does not work:

```
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2 #check convergents c 0
```

# **Proofs**

# The Proof Is Trivial

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Wait, let's do some examples first...

# The Proof Is Trivial

Alright, I am sold!

# **Proving... Please Wait**



# **Proving... Please Wait**



Something seems wrong

# **Now It Is Trivial**

That's better!



<Show two short examples in VS Code>

# **Results**



Definition of (generalized) continued fractions and their evaluation

```
**structure gcf := (head : \alpha) (seq : seq (gcf_pair \rightarrow \alpha))

2 def cf := {g : gcf \alpha // \forall (n : \mathbb{N}) (a : \alpha),

\rightarrow (partial_numerators g).nth n = some a \rightarrow a = 1}

3 def convergents (g : gcf \alpha) (n : \mathbb{N}) : \alpha := ...
```

Computable continued fractions for discrete linear ordered floor fields

Computable continued fractions for discrete linear ordered floor fields

Also works for  $\mathbb{R}$  – just not computable...

# Termination proof for archimedian fields

```
theorem termination_iff_rat [archimedean \alpha] (v : \alpha) \hookrightarrow :

Terminates (get_gcf v) \Leftrightarrow \exists (q : \mathbb{Q}), v = (q : \alpha)
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Terminates (get_gcf v) \leftrightarrow \exists (q : \mathbb{Q}), v = (q : \alpha)
```

#### Including a theorem a mathematician would never prove:

```
theorem translate_rat_get_cf {q : Q}
(v_eq_q : v = q) :
((get_gcf q : gcf Q) : gcf a) = get_gcf v :=
```

#### Finite correctness of the computation

```
theorem get_gcf_finite_correctness
(terminates: Terminates (get_gcf v)):
    ∃ (n : N), v = convergents (get_gcf v) n
```

### Some interesting inequalities, and finally:

```
theorem epsilon_convergence : \forall (\epsilon > (0 : \alpha)), 
2 \exists (N : \mathbb{N}), \forall (n \geq N), 
3 |v - convergents (get_gcf v) n| < \epsilon :=
```

### Some interesting inequalities, and finally:

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```

But sadly no library for sequence limits in Lean:(

**End of the Story** 

 Lean's type system is very expressive and great for definitions...

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- Lean's type system is very expressive and great for definitions...
  - · ...if one knows the gotchas.
- · Support on Zulip is fantastic.
- Existing tactics help a LOT...
  - ...but no integration of automated theorem provers yet.

We Need You!

Help us making interactive theorem proving an even better place!

#### Formalisation can be found at

github.com/kappelmann/lean-continued-fractions

# Formalisation can be found at github.com/kappelmann/lean-continued-fractions

Thanks 
$$+ \frac{1}{for + \frac{1}{your + \frac{1}{attention}}}$$

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Thanks + 
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**Any questions?** 

# Image Sources i

- Salt shaker: Modified from bit.ly/2K8Jw8s
- Link 1: bit.ly/2wMGOwE
- Link 2: bit.ly/2RaypfX
- Link 3: bit.ly/2MNGUPt
- Clock: bit.ly/2HOc9GC
- Melting clock: bit.ly/2MKWknv